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# Trust, Privacy, and Security in Digital Business

Third International Conference, TrustBus 2006 Kraków, Poland, September 2006 Proceedings



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#### Preface

This book presents the proceedings of the Third International Conference on Trust, Privacy and Security in Digital Business (TrustBus 2006), held in Kraków, Połand, September 5-7, 2006. The conference continues from previous events held in Zaragoza (2004) and Copenhagen (2005), and maintains the aim of bringing together academic researchers and industry developers to discuss the state of the art in technology for establishing trust, privacy and security in digital business. We thank the attendees for coming to Kraków to participate and debate the new emerging advances in this area.

The conference programme included two keynote presentations, one panel session and eight technical papers sessions. The keynote speeches were delivered by Jeremy Ward from Symantec EMEA on the topic of "Building the Information Assurance Community of Purpose", and by Günter Karjoth from IBM Research - Zurich, with a talk entitled "Privacy Practices and Economics — From Privacy Policies to Privacy SLAs."

The subject of the panel discussion was "Is Security Without Trust Feasible?" chaired by Leszek T. Lilien from Western Michigan University, USA. The reviewed paper sessions covered a broad range of topics, from access control models to security and risk management, and from privacy and identity management to security protocols. The conference attracted 70 submissions, each of which was assigned to four referees for review. The Programme Committee ultimately accepted 24 papers for inclusion, which were revised based upon comments from their reviews.

We would like to express our thanks to the various people who assisted us in organizing the event and formulating the programme. We are very grateful to the Programme Committee members, and external reviewers, for their timely and rigorous reviews of the papers. Thanks are also due to the DEXA Organizing Committee for supporting our event, and in particular to Mrs. Gabriela Wagner for her help with the administrative aspects. We would also like to thank Sokratis Katsikas, Javier López and Günther Pernul for their past efforts in establishing the conference series, and their valuable advice and assistance in enabling us to take it forward.

Finally we would like to thank all of the authors who submitted papers for the event, and contributed to an interesting set of conference proceedings.

September 2006 Kraków, Poland Simone Fischer-Hübner, Karlstad University, Sweden Steven Furnell, University of Plymouth, UK Costas Lambrinoudakis, University of the Aegean, Greece

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# Towards a UML 2.0 Extension for the Modeling of Security Requirements in Business Processes

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Abstract. Security is a crucial issue for business performance, but usually, it is considered after the business processes definition. Many security requirements can be expressed at the business process level. A business process model is important for software developers, since they can capture from it the necessary requirements for software design and creation. Besides, business process modeling is the center for conducting and improving how the business is operated. This paper contains a description of our UML 2.0 extension for modeling secure business process through activity diagrams. We will apply this approach to a typical health-care business process.

## 1 Introduction

The new business scene, where there are many participants and an intensive use of communications and information technologies, implies that enterprises not only expand their businesses but also increase their vulnerability. As a consequence, with the increase of the number of attacks on systems, it is highly probable that sooner or later an intrusion can be successful [19]. This security violation causes losses. For this reason, it is necessary to protect computers and their systems in the best possible way. Best possible security does not necessarily mean absolute security, but a reasonable high security level in relation to the given limitations [25].

On the other hand, business processes are key to maintain competitiveness. Since, they are the ability of an enterprise to describe, standardize, and adapt the way it reacts to certain types of business events, and how it interacts with suppliers, partners, competitors, and customers [21].

Regardless of the importance of the security notion for companies, this is often neglected in business process models, which usually concentrate on modeling the process in a way that functional correctness can be shown [2] mainly due to the fact that the expert in the business process domain is not an expert in security [9]. Typically, security is considered after the definition of the system. This approach often leads to problems, which most of the times are translated into security

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vulnerabilities [17], which clearly justify the need of increasing the effort in the predevelopment phases, where fixing the bugs is cheaper [14].

If we consider that empirical studies show that it is common at the business process level that customers and end users are able to express their security needs [14], then it is possible to capture at a high level, security requirements easily identifiable by those who models business processes. Besides, requirements specification usually results in a specification of the software system which should be as exact as possible [1], since, effective business process models facilitate discussions among different stakeholders in the business, allowing them to agree on the key fundamentals and to work towards common goals [5].

For business process modeling, there are several languages and notations [8], however, UML (Unified Modeling Language) is a widely accepted standard notation. The most important change of UML 2.0 version with respect to the previous ones has been that of the activity diagrams which improve the business process representation. Our work considers a UML 2.0 extension that allows us to incorporate security requirements into activity diagrams from the perspective of the business analyst. We have considered the security requirements identified in the taxonomy proposed in [7].

Our proposal is based on the MDA (Model Driven Architecture) approach. We will define early requirements identification using UML and this will make it possible to perform independent specifications of the implementation. Moreover, we believe that it is possible to have two different perspectives about security requirements at a high level of abstraction. One of them related to business analysts and the other associated with security experts. In this paper we have deepened in the first perspective.

The structure of the rest of the paper is the following: in Section 2, we will summarize the main issues about security in business processes. In Section 3, we will present a brief overview of UML 2.0 activity diagrams and extensions. In Section 4, we will propose a UML 2.0 extension to represent security requirements. Finally, in Section 5, we will present an example and in Section 6 our conclusion will be drawn.

# 2 Security in Business Process

In spite of the importance of security for business processes, we have found out two problems. The first one is that modeling has not been adequate since, generally, those who specify security requirements are requirements engineers that have accidentally tended to use architecture specific restrictions instead of security requirements [6]. And in the second place, security has been integrated into an application in an ad-hoc manner, often during the actual implementation process [2], during the system administration phase [13] or it has been considered like outsourcing [16].

An approach to model security considering several perspectives is presented in [9]. Authors take into consideration the following perspectives: *static*, about the processed information security, *functional*, from the viewpoint of the system processes, *dynamic*, about the security requirements from the life cycle of the objects involved in the business process, *organizational*, used to relate responsibilities to acting parties within the business process and the *business processes* perspective, that provides us with an integrated view of all perspectives with a high degree of abstraction. Moreover, capturing the security requirements of a system is a hard task that must be established at the initial stages of system development, and business spruces offer a

view of business structure that is very suitable as a basis for the elicitation and specification of security requirements. Business process representations may in this way present in all stages of system development different levels of abstraction appropriate for each stage 1141. Consequently, we believe that business analysts can integrate their view about business security into the business process perspective.

On the other hand, functional security requirements tend to vary depending on the kind of application. This cannot be said about security requirements since any application at the highest level of abstraction will tend to have the same basic kinds of valuable and potentially vulnerable assets [7].

The research works related to security specifications carried out by business domain experts are; (i) scarce [2, 9, 15], (ii) oriented to transaction security [20], (iii) directly oriented to information systems in general [23] or (iv) thought for security and software engineers [16]. Moreover, several works [10, 13, 14, 24] have used UML to perform the specification of security requirements. In these works, activity diagrams have not been used to capture security requirements. However, we believe that it is possible that business analysts can express their security requirements through activity diagrams.

## 3 UML 2.0 Activity Diagrams and UML 2.0 Extensions

UML 2.0 is divided into structural and behavioral specifications. Behavior models specify how the structural aspects of a system change over time. UML has three behavior models: activities, state machines, and interactions. Activities focus on the sequence, conditions, and inputs and outputs for invoking other behaviors, state machines show how events cause changes of object state and invoke other behaviors, and interactions describe message-passing between objects that causes invocation of other behaviors [4].

Activity diagrams are the UML 2.0 elements used to represent business processes and workflows [11]. In UML previous versions, expressivity was limited and this fact confused users that did not use the orientation to objects as an approach for modeling. Now, it is possible to support flow modeling across a wide variety of domains [3]. An activity specifies the coordination of executions of subordinate behaviors, using a control and data flow model. Activities may form invocation hierarchies invoking other activities, ultimately resolving to individual actions [18]. The graphical notation of an activity is a combination of nodes and connectors that allow us to form a complete flow.

On the other hand, the Profiles package contains mechanisms that allow metaclasses from existing meta-models to be extended to adapt them for different purposes. The profiles mechanism is consistent with the OMG Meta Object Facility (MOF) [18]. UML profiles consist of Stereotypes, Constraints and Tagged Values. A stereotype is a model element defined by its name and by the base class to which it is assigned. Constraints are applied to the stereotype with the purpose of indicating limitations (e.g. pre or post conditions, invariants). They can be expressed in natural language, programming language or through OCL (Object Constraint Language). Tagged values are additional meta-attributes assigned to a stereotype, specified as name-value pairs.

Research works related to UML 2.0 profiles and business processes refer to aspects of the business such as Customer, kind of Business Process, Goal, Deliverable and Measure [12], Data Warehouse and its relation to business process dynamic structures [22] or they add semantics to the activities considering organizational aspects that allow us to express resource restrictions during the execution of an activity [11].

# 4 UML 2.0 Extension for Modeling Business Process with Security Requirement

Our proposal allows business analysts to specify security requirements in the business process by using activity diagrams. It is the first part of a security requirements specification that will have later to be complemented by a security analyst. Both perspectives let us enrich the security requirements specifications in business processes.

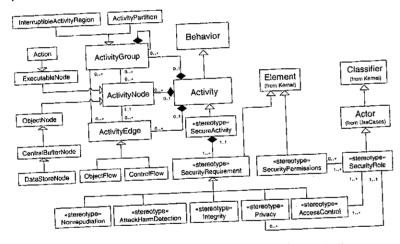


Fig. 1. Extending the UML 2.0 meta-model with security stereotypes

Figure 1 shows the UML 2.0 meta-model extended with stereotypes (in dark) for Secure Activity specifications. A Secure Activity is a stereotype derived from Activity. «SecureActivity» is strongly associated with security requirements stereotypes. «SecurityRequirement» has a composition relationship with «SecureActivity». The proposed notation for «SecurityRequirement» must be complemented by adding it letters that will allow us to identify the type of requirement that is specified.

The stereotypes derived from "SecurityRequirement" can be added to activity diagrams elements. Any security requirement (NR, AD, I, P of AC) can be added to activity diagram elements (see Table 1). For example, an elements requirement can be specified over data store, control flow or object flow.

«SecurityRole» and «SecurityPermissions» are related in different ways, because both can be *obtained* from the UML 2.0 element of activity diagrams (see Table 1). For example, «SecurityRole» can be obtained from activities, partitions or regions specifications, but it is not specified in an explicit way over these activity diagrams elements. «SecurityPermission» is a special case, because, permissions depending on each activity diagram element which they are related to. For example, for Actions object, Execution or CheckExecution operations must be specified (see Table 3).

Table 1. Security Requirements and Activity Diagram Elements

	UML 2.0 element for containment in activity diagrams					
Stereotypes for secure activity specification	Activity	Activity Partition	Interruptible Activity Region	Action	Data StoreNode	Object Flow
Nonrepudiation (NR)	l l					<b>√</b>
AttackHarmDetection(AD)	7	~	~	· /	~	<b>✓</b>
Integrity (I)		1			V	7
Privacy (P)		7				
AccessControl (AC)	<b>1</b>					
Security Role	-	1	-			
SecurityPermissions				~	·	

In addition, we need the definitions of some new data types to be used in tagged value definitions. In Table 2, we will show the new data type stereotypes definitions. All new data types have been derived from the Enumeration Class.

Table 2. New data types

Name	Description	Values associated
SecReqType	It represents a type of security requirement. It must be specified for Non Repudiation, Attack/Harm Detection, Integrity, Privacy or Access Control.	NR, AD, I. P, AC
PerOperations	It is an enumeration for possible operations over objects in activity diagrams. These operations are related to permissions granted over the object	Execution, CheckExecution, Update, Create, Read, Delete, SendReceive, CheckSendReceive
ProtectDegree	It is an abstract level that represents criticality. This degree can be low (I), medium (m) or high (h):	J, m, h
PrivacyType	It consists of anonymity (a) or confidentiality (c).	a, c
AuditingValues	It represents different security events related to the security requirement specification in business processes. They will be used in later auditing	ElementName, SourceName, DestinationName, DateTimeSend, DateTimeReceive, Date, Time, RoleName

Next tables will show the stereotypes for secure activity specifications extensively. Each stereotype specification contains: name, base class, description, notation, constrains and tagged values.

Table 3. Security activity and security requirement stereotypes

Togged Values   Programment	Name: Scenre Act	tivity Description: A secure activity contains security specification related to requirements, role identifications and permissions				
Constrains   Context   SecurityRequirement ->size() >= 1	Base Class: Activ					
SecurityRole   Base Class: Actor (from UseCasss)   Gom access coming and/or privacy specifications   Gom access coming and/or privacy specification and/or interruptible AntivityRole   Gomes   Gome		ti must be associated at least with the				
SecurityRole   Base Class: Actor (from UseCasss)   Gom access coming and/or privacy specifications   Gom access coming and/or privacy specification and/or interruptible AntivityRole   Gomes   Gome	-	strains context SecureActivity inv: self-role specifications. This roles must be obtained				
The role in the security role sereotype can be derived from: Activity, ActivityPartition and/or InterruptibleActivityRegion (see Table 1)	Name: SecurityR	tole Description: to contact the appropriate to the property of the property o				
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Tagged Values   Tagged Value	Base Class	Security Requirement				
Constraints   It can be only specified in the diagram elements indicated in Table 1.		It establishes the need to avoid the demai of any aspect of the				
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Fagged Values  self.AvNr="ElementName" or self.AvNr="SourcoName" or self.AvNr="DestinationName" or self.NvNr="DateTimeGend" or		AvNr: Auditing Values				
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		Sell AVMI - Decelland				

Table 4. Stereotypes specifications for security requirements

Name Base Class	AttackHarmDetection SecurityRequirement	Notation
Description	It indicates the degree to which the attempt or success of attacks or damages is detected, registered and notified. An auditing requirement can be indicated in Comment	AĐ
Constrains	It can be only specified in the diagram elements indicated in Table 1.	
Tagged Values	AvAD: Auditing Values context AttackHaarmDetection inv: self.AvAD="ElementName" or self.AvAD="Date self.AvAD="Time"	or or
Name	fitegrity 2 1 See 1 Was 1 See	Notation
Base Class	SecurityRequirement	. 101211011
Description	It establishes the degree of protection of intentional and non authorized corruption. The elements are protected from intentional corruption. An auditing requirement can be indicated in Comment.	
Constrains	It can be only specified in the diagram elements indicated in Table 1.  The Protection Degree must be specified by adding a lower case letter according to value.	o PDI tagged
Tagged Values	PDI: ProtectDegree Avt: Auddring Values context Integrity inv: self.AvI="Date" or self.AvI="Date" or self.Av	T="Time"
Name.	self.AvI="ElementName" or self.AvI="Date" or self.Av	Notation
Base Class	SecurityRequirement	MALIGN
Description	It indicates the degree to which non authorized parts are avoided to obtain sensitive information. An auditing requirement can be indicated in Comment.	Px
Constrains	* It can be only specified in the diagram elements indicated in Table I.  * A privacy requirement has one security role specification  context Privacy inv: self.SecurityRole -> size() = 1  * The Privacy Type must be specified adding a lower case letter according to Pv tag  privacy type is not specified then annoymity and confidentiality are considered.	ged value. If
Tagged Values	Pv: PrivacyType AvPv: AuditingValues context Privacy inv: self.AvPv="RolcName" or self.AvPv="Date" or self.AvP	v="Time"
Name	AccessControl	Notation
Base Class	SccurityRequirement	
Description	It establishes the need to define and/or intensify the access control mechanisms (identification, authentication and authorization) to restrict access to certain components in an activity diagram. An auditing requirement can be indicated in Comment.	AC
Constrains	R can be only specified in the diagram elements indicated in Table 1.  It is valid only if it is specified at least one security role.  Context AccessControl inv: scif. SecurityRole -> size(	) >= 1
Tagged Values	AvAC: Auditing Values contact AccessControl inv: solf.AvAC="RoleMame" or self.AvAC="Date" or self.AvAC="	Time"

## 5 Example

Our illustrative example (see Figure 2) describes a typical business process for the admission of patients in a health-care institution. In this case, the business analyst identified the following Activity Partitions: Patient, Administration Area (which is a top partition that is divided into Admission and Accounting middle partitions), and the Medical Area (divided into Medical Evaluation and Exams).

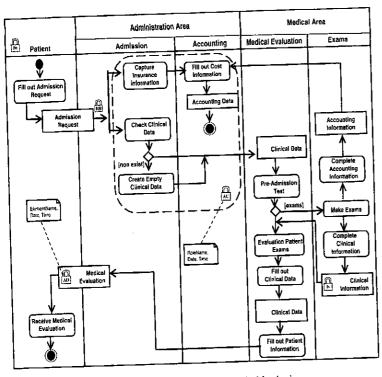


Fig. 2. Admission of Patients in a Medical Institution

The business analyst has considered several aspects of security. He/she has specified «Privacy» (confidentiality) for Activity Partition "Patient", with the aim of preventing the disclosure of sensitive information about Patients. «Nonrepudiation» has been defined over the control flow that goes from the action "Fill Admission Request" to the actions "Capture Insurance Information" and "Check Clinical Data" withthe aim of avoiding the denial of the "Admission Request" reception. «AccessControl» has been defined over the Interruptible Activity Region. A «SecurityRole» can be derived from this specification. Admission/Accounting will be a role. All objects in an interruptible region must be considered for permissions specification (see Table 5). Access control specification has been complemented with audit requirement. This implies that it must register role name, date and time of all events related to the region interruptible. Integrity (high) requirement has specified for Data Store "Clinical Information". Finally, the business analyst has specified Attack Harm Detection with auditing requirement. All events related to attempt or success of attacks or damages are registered (names in this case are clinical information, date and time).

Table 5. «SecurityRole» and «SecurityPermission» specifications

		l'ermissions	
Rolc		Objects	Operations
Admission/Accounting	Action	Capture Insurance Information Fill out Cost information Check Clinical Data Create Empty Clinical Data	Execution CheckExecution Execution Execution
	DataStoreNode	Accounting Data	Update

#### 6 Conclusions and Ongoing Work

The UML 2.0 version, particularly improved for business process representation through activity diagrams, opens an opportunity to incorporate security requirements that allow us to increase this aspect of the systems from early stages in software development. In this paper, we have presented a UML 2.0 extension that allows us to incorporate security requirements into activity diagrams that will increase the scope of the expressive ability of business analysts.

The next step should be that of applying an MDA approach to transform the model (including the security requirements) into most concrete models (i.e. execution models). Therefore, future work must be oriented to enrich the security requirements specifications, improving the UML extension specification to complement it with Well-Formedness Rules and OCL.

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